

# Study of Geometrical structure of Perfect Difference Network (PDN)

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**Abstract:** Geometrical Structure of PDN architecture can be considered as a parameter for critical study for connectivity and complexity of interconnection network. This paper illustrates the new pattern and properties for improving performance and latency time and is also useful for path selection, fast and secure communication system between processors. In this paper the mapping and message passing problem/method is also discussed, which allows the use of standard tools for solving it/problems. Mapping between processor is done for a study of geometrical structure between processor. The main objective of this paper is to simplify the design and increase the performance of Perfect Difference Network. Out of which we have obtained several pattern and property which is mention in this paper.

**Keywords:** Perfect Difference Network, Interconnection Networks, Pattern, message passing, Topological properties, modelling, Bit Manipulation.

## I. INTRODUCTION

In this paper we have studied to explore Interconnection Network to enhance the performance of Perfect Difference Network. Research in Interconnection Network is a broad field, where most results have been achieved for a multitude of Perfect Difference Network (PDN) particularly in parallel and distributed system. A large number of models, algorithms, structure and lemmas have been proposed in the last few years for Perfect Difference Network (PDN).

Perfect difference sets were discussed in 1938 by J.Singer [1] The formulation was in terms of points and lines in a finite projective plane. The perfect difference sets considered for being developed into a interconnect network mainly through works of Behrooz Parhami and Mikhail A Rakov [2]. In their, perfect difference network (PDN) interconnection, they have shown that PDN interconnection scheme is optimal in the sense that it can accommodate an asymptotically maximal number of nodes with smallest possible node degree under the constraint of the network diameter being two. They have compared PDNs and some of their derivatives to interconnection networks with similar cost and performance, including certain generalized hypercube and their hierarchical variants.

Topological properties [3] of perfect difference network compared with the corresponding properties of hypercube. In this scheme, sparse linear system was implemented. It was proved that access function or routing function to map data on hypercube contains topological properties. [3]

Perfect difference networks are a robust high-performance interconnection network for parallel and distributed systems. A comparative study of hypercube and perfect difference network was done, based on topological properties.[4]

An open source NS2 simulator used by Priyanka Wankar et. all [5] to measure the performance analysis and implementation of PDN. Formula  $n=\delta^2+ \delta+1$  is tested on simulator. It was found that simulated PDN and the Mathematical underpinnings makes, it desirable as robust and high-performance.

## II. METHODOLOGY

Diameter can be considered as an important parameter for study of the structural relationship [6] and complexity (time and cost) of interconnection network. Diameter is inversely proportional to the cost of interconnection network but directly proportional to time and meanwhile it affects the communication of interconnection network.

Mapping illustrates how to assign task to a processor of PDN such that it maximize utilization of system resources while minimizing the latency time/communication time. An important point to take into consideration is that the execution of multiple instruction/data on the same processor is not allowed, as this provides a fall of performance in this type of architecture.

## III. RESULT AND DISCUSSION

This section helps to take mapping decision [7] at statically or dynamically by different methods for Perfect Difference Network. Topological Properties are used to map the relationship [7] [8] between processor for Perfect Difference Network. Communication links are bidirectional; it is used for both read and executes program instructions

### • MAPPING AND TRANSFORMATION BETWEEN PROCESSOR OF PERFECT DIFFERENCE NETWORK

In reference of definition of Perfect Difference Network [9] there are  $\delta^2 + \delta + 1$  node/core/processor. Message

passing mechanism [10] between processors of perfect difference network is illustrated below:

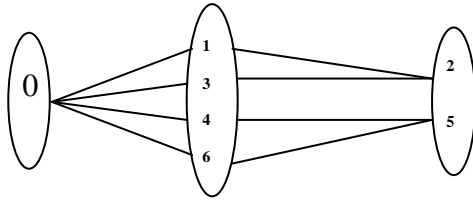


Fig 1: Relationship of  $p_0$  with  $\delta^2 + \delta$  processors

Processor  $p_0$  has the capability to broadcast same instruction for  $p_1, p_3, p_4,$  and  $p_6$  simultaneously.  $p_0$  may use  $P_1$  or  $p_3$  to communicate with  $p_2$  and  $p_4$  or  $p_6$  are used to communicate with  $p_5$ .

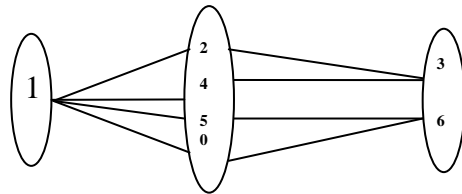


Fig 2: Relationship of  $p_1$  with  $\delta^2 + \delta$  processors

Processor  $p_1$  has the capability to broadcast same instruction for  $p_2, p_4, p_5,$  and  $p_0$  simultaneously.  $P_1$  may use  $P_2$  or  $p_4$  to communicate with  $p_3$  and  $p_5$  or  $p_0$  are used to communicate with  $p_6$ .

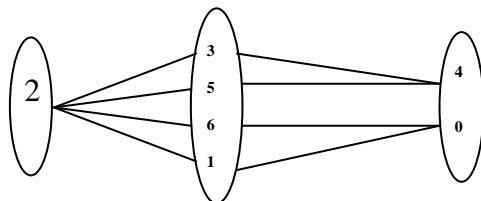


Fig 3: Relationship of  $p_2$  with  $\delta^2 + \delta$  processors

Processor  $p_2$  has the capability to broadcast same instruction for  $p_3, p_5, p_6,$  and  $p_1$  simultaneously.  $P_2$  may use  $P_3$  or  $p_5$  to communicate with  $p_4$  and  $p_6$  or  $p_1$  are used to communicate with  $p_0$ .

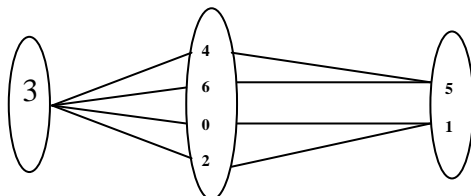


Fig 4: Relationship of  $p_3$  with  $\delta^2 + \delta$  processors

Processor  $p_3$  has the capability to broadcast same instruction for  $p_4, p_6, p_0,$  and  $p_2$  simultaneously.  $P_3$  may use  $P_4$  or  $p_6$  to communicate with  $p_5$  and  $p_0$  or  $p_2$  are used to communicate with  $p_1$ .

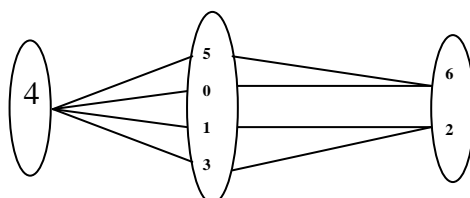


Fig 5: Relationship of  $p_4$  with  $\delta^2 + \delta$  processors

Processor  $p_4$  has the capability to broadcast same instruction for  $p_5, p_0, p_1,$  and  $p_3$  simultaneously.  $P_4$  may use  $P_5$  or  $p_0$  to communicate with  $p_6$  and  $p_1$  or  $p_3$  are used to communicate with  $p_2$ .

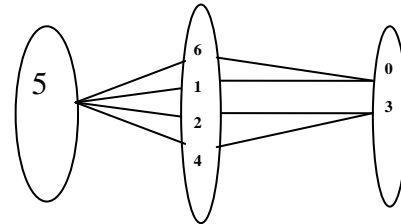


Fig 6: Relationship of  $p_5$  with  $\delta^2 + \delta$  processors

Processor  $p_5$  has the capability to broadcast same instruction for  $p_6, p_1, p_2,$  and  $p_4$  simultaneously.  $P_5$  may use  $P_6$  or  $p_1$  to communicate with  $p_0$  and  $p_2$  or  $p_4$  are used to communicate with  $p_3$ .

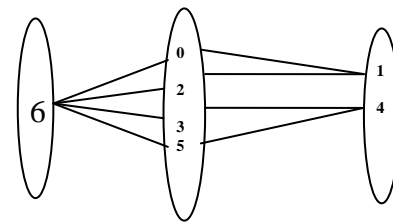


Fig 7: Relationship of  $p_6$  with  $\delta^2 + \delta$  processors

Processor  $p_6$  has the capability to broadcast same instruction for  $p_0, p_2, p_3,$  and  $p_5$  simultaneously.  $P_6$  may use  $P_0$  or  $p_2$  to communicate with  $p_1$  and  $p_3$  or  $p_5$  are used to communicate with  $p_4$ .

These analysis/study are an abstraction of a message passing model to communicate among processors of Perfect Difference Network.

In this analysis it can be concluded that the  $p_i$  has several transitive relations, which are given in the table1.

Table 1: Transitive relation of  $p_i$

| Node/processor/core[ $p_i$ ] | Node having transitive relation |
|------------------------------|---------------------------------|
| 0                            | 2,5                             |
| 1                            | 3,6                             |
| 2                            | 4,0                             |
| 3                            | 5,1                             |
| 4                            | 6,2                             |
| 5                            | 0,3                             |
| 6                            | 1,4                             |

Table 2: PDS relationship between  $p_i$  and  $p_{i_{transitive}}$ .

| PDS( $p_i$ ) | PDS( $p_{i_{transitive}}$ ) |
|--------------|-----------------------------|
| (0-0)        | (3-1) , (1-3)               |
| (1-0)        | (3-0) , (0-1)               |
| (3-1)        | (0-3) , (0-0)               |
| (3-0)        | (1-3) , (1-0)               |
| (0-3)        | (0-1) , (3-1)               |
| (1-3)        | (0-0) , (3-0)               |
| (0-1)        | (1-0) , (0-3)               |

Algorithm 1: **MACBPPDN**

**BEGIN**

STEP 1: Declare all the variable

STEP 2: Assign memory allocation to the PDS.

STEP 3: Input value of  $\delta$ .

STEP 4: Validate the  $\delta$  value

- (i) Is value of  $\delta$  is prime or power of prime.
- (ii) The PDS value of particular delta ( $\delta$ ) is available or not.

STEP 5: Input Source ( $p_s$ ) and Destination ( $p_d$ ).

STEP 6: Evaluate the formula/ equation

// using structural relation study

STEP 7: Display the Result.

**END**

The main objective of this algorithm is to select paths of small total delay for each data/instruction. If there were no queuing delays along any link, the path selection would be simple. The time for a data/instruction to travel between two processors is  $O(2)$ , assuming no queuing delays at the links.

• **BOOLEAN OPERATION ON STATE OF PDN**

While doing mapping several structure [11][12] and pattern have been find out. Bit representation/state of processors are useful for designing circuits [13] and it may model a logical system with a single equation. The three operation AND, OR, EOR (exclusive or) are used in the state of processors which is found by transitive table [7][14] of Perfect Difference Network. The main objective of this section is to simplify the design and increase the performance of Perfect Difference Network. Out of which we have obtained several pattern and property which is mention below.

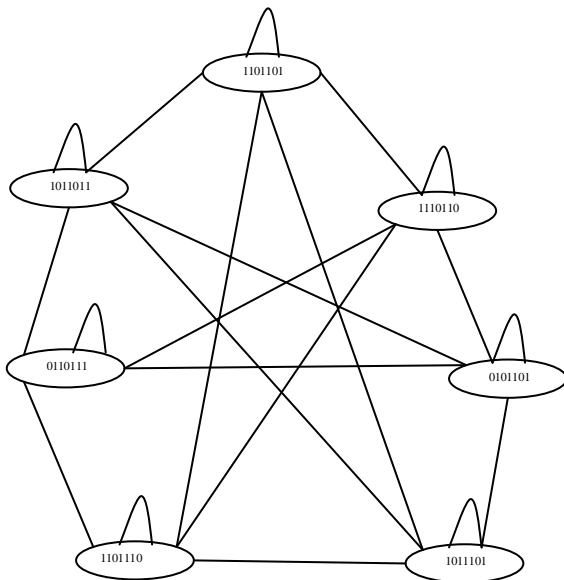


Fig 8: State Diagram

Table 3: AND operation on fstate(PDN(TT))

| AND | 1 | 1 | 0 | 1 | 1 | 0 | 1 |
|-----|---|---|---|---|---|---|---|
| 1   | 1 | 1 | 0 | 1 | 1 | 0 | 1 |
| 1   | 1 | 1 | 0 | 1 | 1 | 0 | 1 |
| 0   | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1   | 1 | 1 | 0 | 1 | 1 | 0 | 1 |
| 1   | 1 | 1 | 0 | 1 | 1 | 0 | 1 |
| 0   | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1   | 1 | 1 | 0 | 1 | 1 | 0 | 1 |

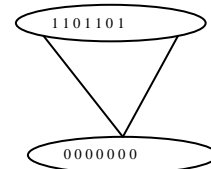


fig 9: relationship between state of table 3.

Fig. 9 shows the relationship between state of AND operation, there are two bits (0) is match (3<sup>rd</sup> & 6<sup>th</sup> position of state) counting starts from left to right.

Table 4: OR operation on fstate(PDN(TT))

| OR | 1 | 1 | 0 | 1 | 1 | 0 | 1 |
|----|---|---|---|---|---|---|---|
| 1  | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| 1  | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| 0  | 1 | 1 | 0 | 1 | 1 | 0 | 1 |
| 1  | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| 1  | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| 0  | 1 | 1 | 0 | 1 | 1 | 0 | 1 |
| 1  | 1 | 1 | 1 | 1 | 1 | 1 | 1 |

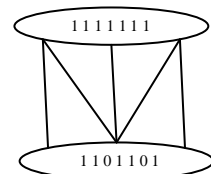


fig 10: relationship between state of table 4.

Fig.10 shows the relationship between state of OR operation, there are five bits (1) is match (0<sup>th</sup>, 1, 3<sup>rd</sup>, 4<sup>th</sup>, & 6<sup>th</sup> position of state) counting starts from left to right.

Table 5: X-OR operation on fstate(PDN(TT))

| E-OR | 1 | 1 | 0 | 1 | 1 | 0 | 1 |
|------|---|---|---|---|---|---|---|
| 1    | 0 | 0 | 1 | 0 | 0 | 1 | 0 |
| 1    | 0 | 0 | 1 | 0 | 0 | 1 | 0 |
| 0    | 1 | 1 | 0 | 1 | 1 | 0 | 1 |
| 1    | 0 | 0 | 1 | 0 | 0 | 1 | 0 |
| 1    | 0 | 0 | 1 | 0 | 0 | 1 | 0 |
| 0    | 1 | 1 | 0 | 1 | 1 | 0 | 1 |
| 1    | 0 | 0 | 1 | 0 | 0 | 1 | 0 |

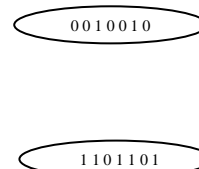


fig 11: relationship between state of table 5.

Note: No node in PDN found to be mutually exclusive which shows that PDN is well connected to each other hence proved that PDN is robust.

From the above studied it can be concluded that, if we apply the EOR operation between the state of OR, we get an indirect relation of processor  $p_i$  and if we apply the EOR operation between the state of AND, we get a direct relation of processor  $p_i$ .

• INTERCONNECTION WITH NODE X

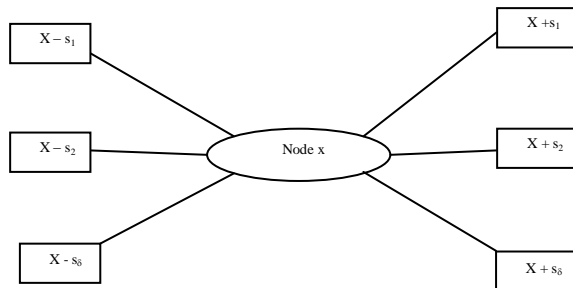


Fig.12: interconnection with  $p_x$ .

This section shows the direct interconnection network [15] with  $n = \delta^2 + \delta + 1$  nodes based on the normal form perfect difference set  $\{s_0, s_1, \dots, s_\delta\}$  with order  $\delta$ .

Node  $x$  is connected via directed links to nodes  $i \pm 1$  and  $i \pm j \pmod n$ , for  $2 \leq j \leq \delta$  [2][3][6]. For each link from node  $x$  to others, the reverse link from others to node  $x$  is also possible. The communication between nodes in a PDN is Symmetric and from the above studied it can be already proved the communication between nodes in a PDN is also transitive [16][17][18].

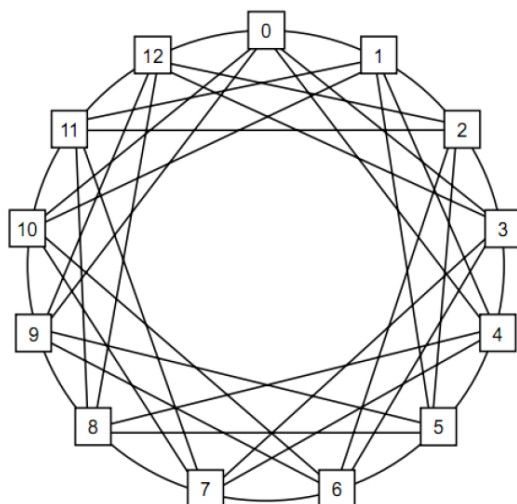


Fig 13: PDN with  $n = 13$  nodes based on the perfect difference set  $\{0, 1, 3, 9\}$ .

IV. CONCLUSION

The focus of this paper is on modelling the mapping and message passing problem which allows the use of standard tools for solving it. Mapping between processor is done for a study of geometrical structure between processor. PDN can be used for Network-on-Chip (NoC). NoC is a technique based on System-on-Chip (SoC) which attempts

to provide high performance nanoscale architectures [18]. Network can be made reusable by separating the communication infrastructure from computing resources.

A perfect difference network is a robust, high performance interconnection network for parallel and distributed computation. PDNs may be desirable for large networks with wired connectivity, but definitely they do offer attractive alternatives for wireless, nanophotonic and optical interconnections and as smaller component networks in hierarchical architectures.

From the above observation/analysis it can be conclude that diameter should have a balanced value. There are  $(\delta^2 + \delta + 1)$  processors having transitive operation with each other and each nodes come in two times in transitive relation (as in table 1).

We have also observed that the  $p_i$  has not always direct relation with its mirror image because the  $p_i$  has direct relation with  $2\delta$  processors and transitive relation with others. The mirror image of  $p_i$  can have its transitive relationship with  $p_i$ . Relational study between node  $i$  and its connected node also has transitive relation using PDS

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